

Transportation

⇒ North West Corner Rule

Q	Factory	^{starting} w_1	w_2	w_3	w_4	Factory Capacity
	F_1	(19) ⁵	30	50	10	(7) ²
	F_2	70	30	40	60	9
	F_3	40	8	70	20	8
	warehouse requirement	(5) ⁰	8	7	14	34

34

Sum should be same otherwise it is an unbalanced problem

write minimum
number on 19

as 5 is over

write 0 below

and 2 is remaining

out of 7 after
removing 5

now w_1 is satisfied so its
over

Factory	w_1	w_2	w_3	
F_1	(30) ²	50	10	(2) ⁰
F_2	30	40	60	9
F_3	8	70	20	18
	(8) ⁶	7	14	

remove
row/column
which becomes
0

	w_1	w_2	w_3	
F_1	(30) ⁶	40	60	(9) ³
F_2	8	70	20	18
	(8) ⁰	7	14	

	w_1	w_2	
F_1	40 ³	60	(3) ⁰
F_2	70	20	18
	(7) ⁴	14	

	w_1	w_2	
F_1	(70) ⁴	20 ¹⁴	(18) ¹⁴
	(4) ⁰	14	

in last table
put remaining
number at the same digit

now in the base table write numbers received by each value

	w_1	w_2	w_3	w_4	
F_1	19 ^③	30 ^②	50	10	7
F_2	70	30 ^⑥	40 ^③	60	9
F_3	40	8	70 ^④	20 ^⑭	18
	5	8	7	14	34

check for $m+n-1$ allocation

$$3 + 4 - 1 = 6$$

so, it is basic feasible solution

$$\begin{aligned}
 \text{Total Cost} &= 19 \times 5 + 30 \times 2 + 30 \times 6 + 40 \times 3 \\
 &\quad + 70 \times 4 + 20 \times 14 \\
 &= 85 + 60 + 180 + 120 + 280 + 280 \\
 &= 1015
 \end{aligned}$$

Least Cost Method

	w_1	w_2	w_3	w_4	S
F_1	19	30	50	10	7
F_2	70	30	40	60	9
F_3	40	⑧	70	20	18 ¹⁰ ←
D	5	⑧	7	14	

circle minimum
delivery cost

out of 8 and 18
min is 8

	w_1	w_3	w_4		w_1	w_3	w_4		
F_1	19	50	10	70	F_2	70	40	60	9
F_2	70	40	60	9	F_3	40	70	20	10 ³
F_3	40	70	20	10		5	7	70	
	5	7	14	7					

$$\begin{array}{cc}
 & w_1 \quad w_3 \\
 F_1 & 70 \quad 40^7 \quad 9^2 \\
 F_3 & 40 \quad 70 \quad 3 \\
 & 5 \quad 7_0
 \end{array}$$

we can allocate on any of the 40

$$\begin{array}{cc}
 & w_1 \\
 F_1 & 70^2 \quad 2 \\
 F_3 & 40^3 \quad 310^5 \\
 & 5^0
 \end{array}$$

on last step we give them their correspondy value

	w_1	w_2	w_3	w_4	
F_1	19	30	50	10	7
F_2	70	30	40	60	9
F_3	40	8	70	20	18
	5	8	7	14	

$$\begin{aligned}
 \text{Total Cost} &= 70 \times 2 + 40 \times 3 + 8 \times 8 + 40 \times 7 + 7 \times 10 \\
 &\quad + 20 \times 7 \\
 &= 314
 \end{aligned}$$

VAM

	w_1	w_2	w_3	w_4	
F_1	19	30	50	10	27
F_2	40	30	40	60	9
F_3	40	8	70	20	18
	5	8	7	14	34

check the lowest

(9) value and 2nd lowest (10) in a column and (12) write their diff.

(21) (22) (10) (10)

then look now

now check for highest penalty

in the column / row

check for minimum

now do just like did in

previous method

but in this case we

will get new penalty each time

	w_1	w_3	w_4	
F_1	19	50	10	7 (9)
F_2	40	40	60	9 (20)
F_3	40	70	20	10 (20)
	5	7	14	

(21) (10) (10)

	w_3	w_4	
F_1	50	(10)	2 (40)
F_2	40	60	9 (20)
F_3	70	(20)	10 (50) ←
	7	14	
	(10)	(10)	

	w_3	w_4	
F_1	50	(10) ²	20 (40)
F_2	40	60	9 (20)
	7	4 ²	
	(10)	(50)	
		↑	

	w_3	w_4
F_2	40 ⁷	60 ²
	7 ⁰	2 ⁰

	w_1	w_2	w_3	w_4	
F_1	195	30	50	10 ²	7
F_2	40	30	40 ⁷	60 ²	9
F_3	40	8 ⁸	70	20 ¹⁰	18
	5	8	7	14	

Total Cost = $19 \times 5 + 8 \times 8 + 40 \times 7 + 10 \times 2 + 60 \times 2 + 10 \times 7$
 $= 779$

Optimality Test.

To check if cost can be further lowered

Modi Method

	D ₁	D ₂	D ₃	D ₄	
O ₁	1 ²⁰	2 ¹⁰	1 ¹⁰	4	30
O ₂	3	3 ²⁰	3 ²⁰	1 ¹⁰	50
O ₃	4	2 ²⁰	5	9	20
	20	40	30	10	

Total Cost = 180

For allocated cells $C_{ij} = \cancel{u_i + v_j} u_i + v_j$

	D ₁	D ₂	D ₃	D ₄	
O ₁	1 ²⁰		1 ¹⁰		$u_1 \Rightarrow -2$
$\Rightarrow$ O ₂		3 ²⁰	3 ²⁰	1 ¹⁰	$u_2 = 0$
O ₃		2 ²⁰			$u_3 \Rightarrow -1$
	v_1	v_2	v_3	v_4	
	3	3	3	1	

These check which row/column has highest no. of allocation and assume its value 0

Sum should be equal to cost

For non-allocated cells $C_{ij} - (u_i + v_j) \geq 0$

	D ₁	D ₂	D ₃	D ₄	
O ₁		+2 ¹	+4 ⁻¹		$u_1 (-2)$
O ₂	3 ³⁺				$u_2 (0)$
O ₃	4 ²⁺	+5 ²⁺	9 ⁰		$u_3 (-1)$
	v_1	v_2	v_3	v_4	
	(3)	(3)	(3)	(1)	

As $C_{ij} - (u_i + v_j) \geq 0$

So, The cost 180 is optimal

Q when non-allocated becomes negative

	w_1	w_2	w_3	w_4	
F_1		30^{-2}	50^0		$u_1 = 10$
F_2	70^{+8}	30^{+8}			$u_2 = 60$
F_3	40^{+8}		70^0		$u_3 = 20$
	v_1	v_2	v_3	v_4	
	9	-12	-20	0	

at (2,2) cell $c_{22} - (u_2 + v_2) \leq 0$

Now allocate 0 at (2,2) cell
it will done in the table of allocated

	19^5	10^2	7	
	$30^0 \leftarrow 40^7$	60^{2-0}	9^{+0}	check above and below
	8^{8+0}	20^{10+0}	18^{-0}	30^0 far remove
	5	8^{+0}	11^{+0}	0 from $3+0$

then go to the row

wherever 0 takes a new position outside constant
will have effect

Now of all the values of the edges take the -ve. for
minimizing

$$\min(8-0, 2-0) = 0$$

$$8=0 \quad 2=0 \leq \text{minimum}$$

$$\text{eg } 0 = 2$$

Now find cost by putting value of 0 in allocated
table.

Unbalanced

Check the sum of Demand and supply
if whichever is less give it a dummy
row/column

	D ₁	D ₂	D ₃	
	4	8	8	76
	16	24	16	82
	8	16	24	77
	72	102	41	

diff = 20

Give dummy column all values 0
and outside value is excess value

	D ₁	D ₂	D ₃	D ₄	
	4	8	8	0	76
	16	24	16	0	82
	8	16	24	0	77
	72	102	41	20	

Now solve normally

Prohibited Routes

Maximization

Whenever it is mentioned transportation cost
minimize it and when it is profit maximize it

Table give for max can't be done we have to convert it for min

For conversion take highest cost and write it's difference with other costs

	D ₁	D ₂	D ₃	
S ₁	8	3	6	120
S ₂	5	10	12	80
S ₃	3	9	10	80
	150	80	50	

	D ₁	D ₂	D ₃	
S ₁	7	10	9	120
S ₂	0	5	3	80
S ₃	12	6	5	80
	150	80	50	

	D ₁	D ₂	D ₃	
S ₁	7 ⁷⁰	10 ⁵⁰	9	120
S ₂	0 ⁸⁰	5	3	80
S ₃	12	6 ³⁰	5 ⁵⁰	80
	150	80	50	

now revert back to original table with allocate at pos.

	D ₁	D ₂	D ₃	
S ₁	8 ⁷⁰	5 ⁵⁰	6	
S ₂	15 ⁸⁰	10	12 ⁵⁰	
S ₃	3	9 ³⁰	10	

Now just multiply:

Alternate Solution

When checking under Modi, if in second phase difference of any cost i.e. $c_{ij} - (u_i + v_j)$ is 0 alternate exists

	D_1	D_2	D_3
S_1	+6 ³⁵	3 ²⁵	5
S_2	-5 ⁴⁰	2 ⁴⁰	2 ⁴⁰
S_3	12	7 ⁸⁵	8

The cell for which we get zero is the starting point, starting ph. will be +ve then next will be -ve

We will take change to be -ve current loop

$$\text{Change} = -3 \quad (\text{as } 3 < 5)$$

Now add change where +ve sign and subtract where -ve sign

	D_1	D_2	D_3	
S_1	9 ³⁵	0 ²⁵	5	65
S_2	2 ⁴⁰	5 ⁴⁰	2 ⁴⁰	0
S_3	12	7 ⁸⁵	8	20
	15	65	40	

take $u=0$
for row having
maximum allocation